

Max-Diversity Affine Precoding for the Noncoherent Doubly Dispersive Channel

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Motivating Questions:

1. What is the maximum achievable diversity order for communication over an *unknown* time/frequency-selective channel?
2. How should the transmitted signal be designed to facilitate maximum diversity reception?

System Model:

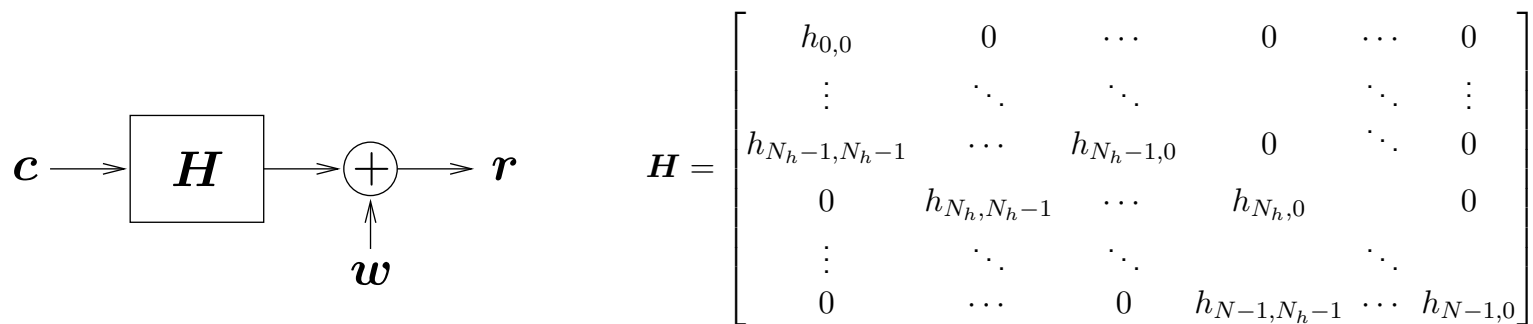
$$r_n = \sum_{l=0}^{N_h-1} h_{n,l} c_{n-l} + w_n$$

$\mathbf{r} := [r_0, \dots, r_{N-1}]^T$: received samples

$\mathbf{c} := [c_0, \dots, c_{N-1}]^T$: coded symbols

$\mathbf{w} := [w_0, \dots, w_{N-1}]^T$: noise samples, $\mathcal{CN}(\mathbf{0}, \sigma^2 \mathbf{I})$

\mathbf{H} : LTV channel matrix



Karhunen-Loève Basis Expansion Model:

KL-BEM of l^{th} -tap trajectory over N -sample block:

$$\mathbf{h}_l := \begin{bmatrix} h_{0,l} \\ \vdots \\ h_{N-1,l} \end{bmatrix} = \mathbf{B}_l \boldsymbol{\theta}_l, \quad \boldsymbol{\theta}_l \in \mathbb{C}^{N_b}$$

N_b : Temporal degrees-of-freedom per tap

WSSUS Rayleigh channel assumption:

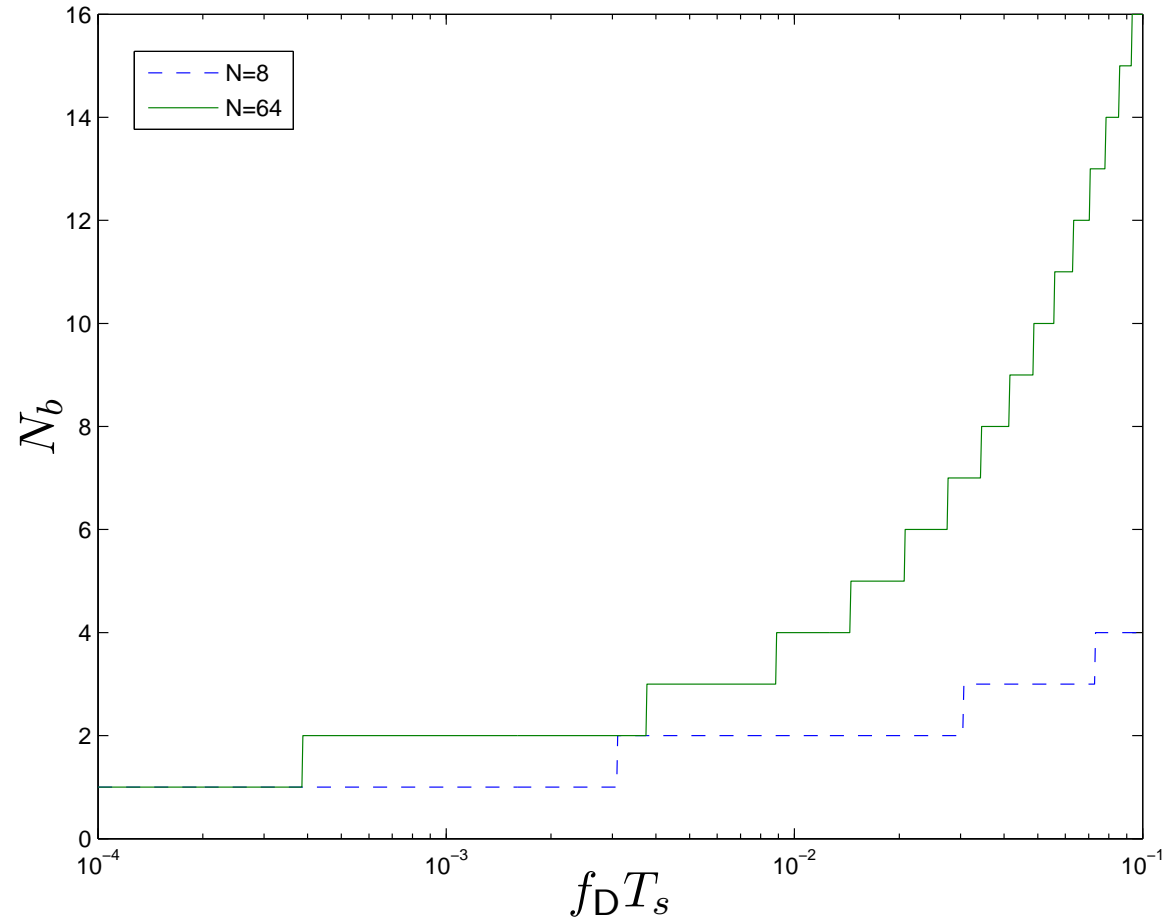
N_h : Number of taps

$$\mathbf{B} = \mathbf{B}_l \quad \forall l \in \{0, \dots, N_h - 1\}$$

$$\boldsymbol{\theta} := \begin{bmatrix} \boldsymbol{\theta}_0 \\ \vdots \\ \boldsymbol{\theta}_{N_h-1} \end{bmatrix} \sim \mathcal{CN}(\mathbf{0}, \mathbf{R}_\theta), \quad \mathbf{R}_\theta \text{ has full rank } N_h N_b$$

Per-tap DoF N_b vs. normalized Doppler spread $f_D T_s$:

of eigenvalues within 30dB of largest eigenvalue (Jakes spectrum):



Notice that $N_b \ll N$.

Noncoherent ML Decoding:

Goal: Estimate $\mathbf{c} \in \mathcal{C}$ from $\mathbf{r} = \mathbf{H}\mathbf{c} + \mathbf{w}$ assuming $\{\mathbf{H}, \mathbf{w}\}$ are unknown but statistics $\{\mathbf{B}, \mathbf{R}_\theta, \sigma^2\}$ are known.

Writing the received vector as

$$\mathbf{r} = \mathbf{C}_B \boldsymbol{\theta} + \mathbf{w},$$

where matrix \mathbf{C}_B is composed of coded symbols \mathbf{c} and basis vectors \mathbf{B} , the noncoherent ML estimate can be written

$$\hat{\mathbf{c}}_{\text{ML}} = \arg \min_{\mathbf{c} \in \mathcal{C}} \mathbf{r}^H \boldsymbol{\Phi} \mathbf{r} - \sigma^2 \log \det(\mathbf{C}_B^H \mathbf{C}_B + \sigma^2 \mathbf{R}_\theta^{-1})$$

\mathcal{C} : set of code vectors

$$\boldsymbol{\Phi} := (\mathbf{C}_B \mathbf{R}_\theta \mathbf{C}_B^H + \sigma^2 \mathbf{I}_N)^{-1}$$

Pair-Wise Error Probability:

Lemma 1 Say $\{\mathbf{C}_B^{(k)}, \mathbf{C}_B^{(l)}\}$ are two possibilities for \mathbf{C}_B . If the matrix

$$\mathbf{M}_{kl} := \mathbf{C}_B^{(k)H} \left(\mathbf{I}_N - \mathbf{C}_B^{(l)} (\mathbf{C}_B^{(l)H} \mathbf{C}_B^{(l)})^{-1} \mathbf{C}_B^{(l)H} \right) \mathbf{C}_B^{(k)}$$

is full rank, then, at high SNR,

$$PWE P_{kl} = \left(\frac{1}{\sigma^2} \right)^{-N_h N_b} \det(\mathbf{R}_\theta \mathbf{M}_{kl})^{-1} \binom{2N_h N_b - 1}{N_h N_b}.$$

Furthermore, \mathbf{M}_{kl} has full rank $N_h N_b$ if and only if $[\mathbf{C}_B^{(k)}, (\mathbf{C}_B^{(l)} - \mathbf{C}_B^{(k)})]$ has full rank $2N_h N_b$.

Main points:

1. $N_h N_b$ is the maximum achievable diversity order.
2. Max-diversity requires full-rank $[\mathbf{C}_B^{(k)}, (\mathbf{C}_B^{(l)} - \mathbf{C}_B^{(k)})] \forall k \neq l$ which requires $N \geq 2N_h N_b$.

Linear Precoding:

Say that the code vectors are generated via

$$\boxed{\mathbf{c} = \mathbf{P}\mathbf{s}} \quad \mathbf{s} \in \mathcal{S} \subset \mathbb{C}^{N_s}$$

where $\mathbf{P} \in \mathbb{C}^{N \times N_s}$ is a linear precoding matrix and \mathbf{s} is a symbol vector.

Lemma 2 *Linear precoding does not facilitate maximum-diversity decoding whenever the symbol vector alphabet \mathcal{S} contains elements that differ by no more than a scale factor (e.g., uncoded QAM or PSK).*

Affine Precoding:

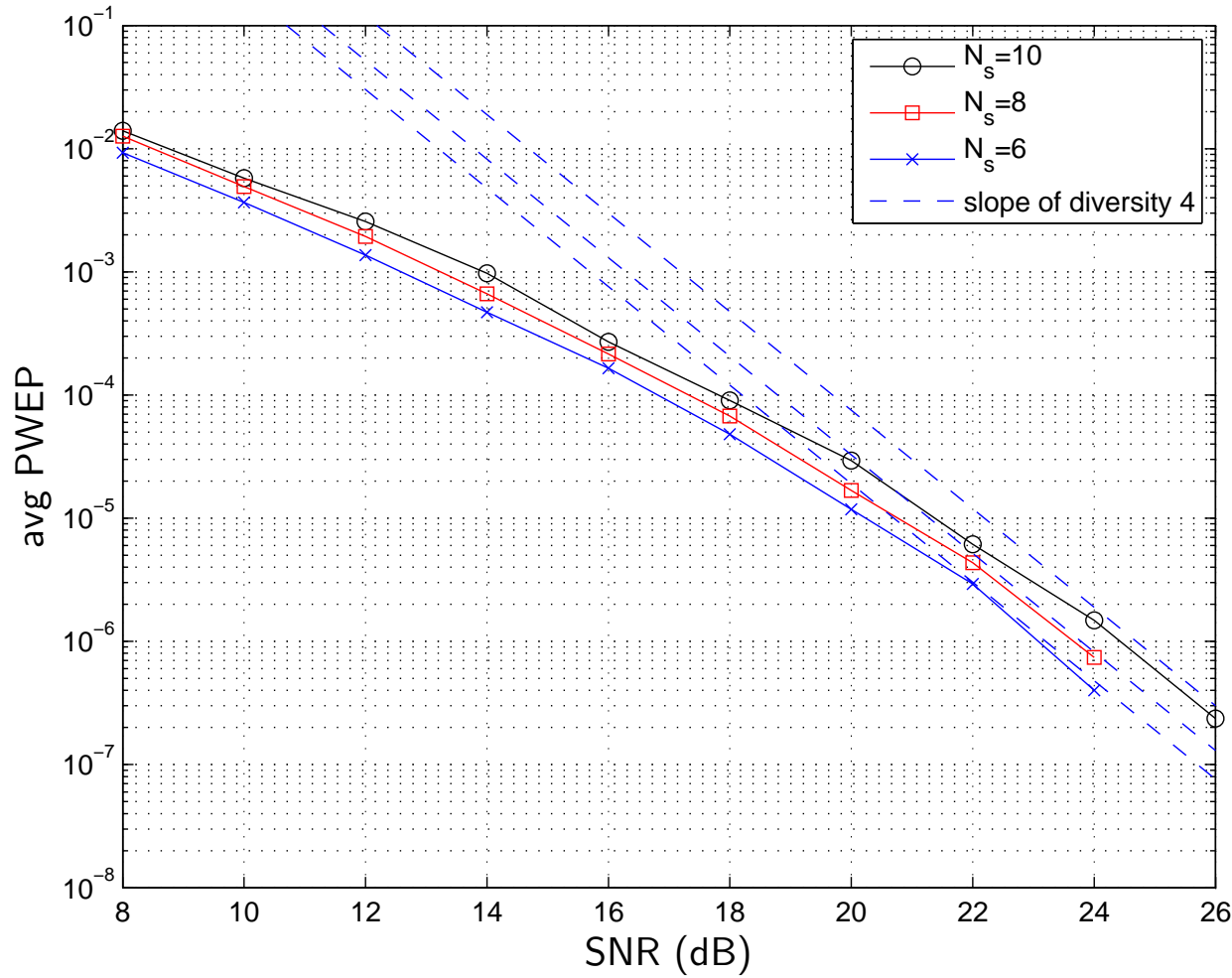
$$\boxed{\mathbf{c} = \mathbf{P}\mathbf{s} + \mathbf{t}} \quad \mathbf{s} \in \mathcal{S} \subset \mathbb{C}^{N_s}.$$

Lemma 3 *If $N \geq 2N_h N_b$ and if the matrix created from the last $N - N_h + 1$ rows of \mathbf{B} is full rank, then choosing $[\mathbf{P}, \mathbf{t}]$ randomly ensures that $[\mathbf{C}_B^{(k)}, (\mathbf{C}_B^{(l)} - \mathbf{C}_B^{(k)})]$ is full-rank w.p.1.*

Main points:

1. *Almost any* affine precoder provides maximum diversity!
2. There are no restrictions on the data rate N_s/N !
3. The rank condition on \mathbf{B} is mild. (It requires that the first $N_h - 1$ samples of the N -sample tap trajectory are non-essential to experiencing the N_b degrees-of-freedom.)

Numerical Example:



$N = 8$

$N_h = 2$

$N_b = 2$

Jakes fading

$\mathcal{S} = \text{BPSK}$

Systematic Affine Precoding:

$$\boxed{\mathbf{c} = \mathbf{P}\mathbf{s} + \mathbf{t}} \quad \text{with} \quad \mathbf{P} = \begin{bmatrix} \mathbf{I}_{N_s} \\ \mathbf{P}' \end{bmatrix}, \quad \mathbf{P}' \in \mathbb{C}^{N_p \times N_s}$$

Lemma 4 *If $N_p \geq N_h N_b - 1$, if $N \geq 2N_h N_b$, and if matrix created from the last $N_p - N_h + 1$ rows of \mathbf{B} is full-rank, then choosing $[\mathbf{P}', \mathbf{t}]$ randomly ensures that $[\mathbf{C}_B^{(k)}, (\mathbf{C}_B^{(l)} - \mathbf{C}_B^{(k)})]$ is full-rank w.p.1.*

Main points:

1. Systematic affine precoding facilitates *fast decoding*.
2. With $N_p \geq N_h N_b - 1$, almost any precoder provides max-diversity!
3. Rate limitation: $\frac{N_s}{N} \leq 1 - \frac{N_h N_b - 1}{N}$.
4. As before, the rank condition on \mathbf{B} is mild.

Conclusions:

For noncoherent communication over a WSSUS time/frequency-selective channel with N_h delay taps and N_b temporal degrees-of-freedom per tap,

1. the maximum diversity order is $N_h N_b$,
2. block lengths $N \geq 2N_h N_b$ facilitate max-diversity reception,
3. linear precoding does *not* facilitate max-diversity reception,
4. *almost any* affine precoder facilitates max-diversity reception *at any rate* $\frac{N_s}{N}$,
5. *systematic* affine precoding facilitates max-diversity at rates $\frac{N_s}{N} \leq 1 - \frac{N_h N_b - 1}{N}$ while simplifying the decoding task.