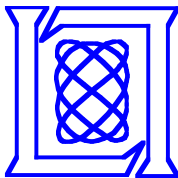


Full-Duplex MIMO Relaying: Achievable Rates Under Limited Dynamic Range

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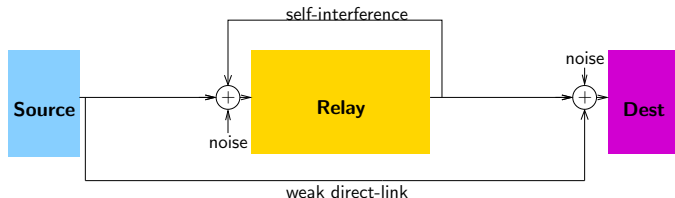
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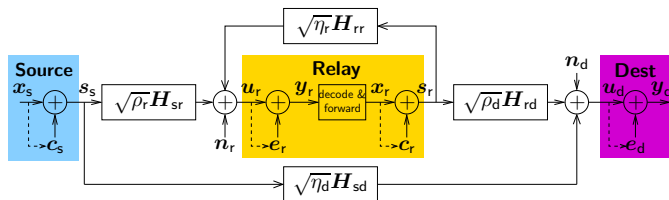
[§] This work was sponsored by the Defense Advanced Research Projects Agency under Air Force contract FA8721-05-C-0002.

Introduction



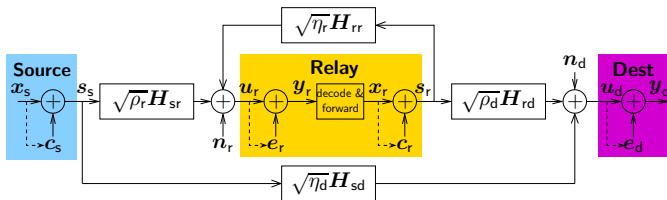
- Source communicates to Destination through **decode and forward Relay**
- **MIMO** at all terminals
- Relay operates in **full-duplex** mode
- Fundamental challenges:
 - **high self-interference** (as high as 100dB!)
 - **limited dynamic range** due to non-ideal transmitter and receiver hardware (power amp noise, non-linearities in ADC/DAC, oscillator phase noise, AGC noise)
- Fundamental question:
 - **What is the maximum achievable rate of such systems?**

System Model



- H_{ij} are MIMO Rayleigh fading propagation channels, assumed to be unknown and static.
- n_i is AWGN thermal noise of unit variance.
- ρ_i represents SNR and η_i represents INR.
- Dynamic range limitation modeled by signal-power dependent additive interference c_j at transmitters and e_i at receivers.
 - Facilitates tractable achievable-rate analysis.
 - Recent work (e.g. Rice, Lincoln) confirms the fidelity of this model.

Dynamic Range Limitation Model



- Each receive chain is corrupted by **additive Gaussian interference** with **power proportional to the intended receive power**; similar for each transmit chain

$$\mathbf{e}_i(t) \sim \mathcal{CN}(\mathbf{0}, \beta \text{diag}(\Phi_i)), \quad \mathbf{e}_i(t) \perp \mathbf{u}_i(t), \quad \mathbf{e}_i(t) \perp \mathbf{e}_i(t') \Big|_{t' \neq t}$$

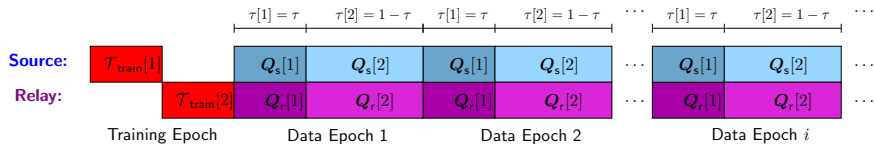
$$\mathbf{c}_j(t) \sim \mathcal{CN}(\mathbf{0}, \kappa \text{diag}(\mathbf{Q}_j)), \quad \mathbf{c}_j(t) \perp \mathbf{x}_j(t), \quad \mathbf{c}_j(t) \perp \mathbf{c}_j(t') \Big|_{t' \neq t}$$

where $\Phi_i = \text{Cov}(\mathbf{u}_i)$ and $\mathbf{Q}_j = \text{Cov}(\mathbf{x}_j)$.

Transmission Protocol

- During Epoch i , the source communicates the i^{th} packet to the relay, while the relay simultaneously communicates the $(i - 1)^{st}$ packet to the destination. \leadsto Enables **full-duplex** communication.
- Before the first data epoch, we have a **training epoch** where we perform **least-squares** channel estimation.
- Data communication parameters (e.g. transmit covariance matrices) are designed to **maximize the achievable rate**.

Two Periods Per Epoch



- We allow two **distinct transmit covariance matrices** per data epoch.
- The two periods per data epoch can **differ in duration**.

Partial Interference Cancellation

- We show that the relay's received signal can be modeled as

$$\mathbf{y}_r(t) = \sqrt{\rho_r} \hat{\mathbf{H}}_{sr} \mathbf{x}_s(t) + \mathbf{v}_r(t)$$

where \mathbf{v}_r is the **aggregate interference** including transmitter/receiver dynamic-range induced self-noise, channel-estimation error, and thermal noise. Similarly, we can write \mathbf{y}_d with interference \mathbf{v}_d .

- We write the relay's aggregate interference as

$$\begin{aligned} \mathbf{v}_r(t) \triangleq & \sqrt{\eta_r} \hat{\mathbf{H}}_{rr} \mathbf{x}_r(t) + \sqrt{\rho_r} \hat{\mathbf{H}}_{sr} \mathbf{c}_s(t) - \mathbf{D}_{sr}^{\frac{1}{2}} \tilde{\mathbf{H}}_{sr} (\mathbf{x}_s(t) + \mathbf{c}_s(t)) + \mathbf{n}_r(t) \\ & + \sqrt{\eta_r} \hat{\mathbf{H}}_{rr} \mathbf{c}_r(t) - \mathbf{D}_{rr}^{\frac{1}{2}} \tilde{\mathbf{H}}_{rr} (\mathbf{x}_r(t) + \mathbf{c}_r(t)) + \mathbf{e}_r(t) \end{aligned}$$

where $\sqrt{\eta_r} \hat{\mathbf{H}}_{rr} \mathbf{x}_r(t)$ is known by the relay and can be eliminated using **interference cancellation**.

Lower-Bounding the Achievable Rate

- Mutual information characterization is complicated by the fact that the **aggregate interference** \mathbf{v}_i is **non-Gaussian** when channel-estimation error is non-zero.
- We therefore **lower-bound** the mutual information by replacing \mathbf{v}_i with a Gaussian noise of identical covariance, i.e.,

$$\underline{I}_{\text{sr}}(\mathcal{Q}[l]) = \log \det (\mathbf{I} + \rho_r \hat{\mathbf{H}}_{\text{sr}} \mathbf{Q}_s[l] \hat{\mathbf{H}}_{\text{sr}}^H \hat{\Sigma}_r^{-1}[l])$$

where $\hat{\Sigma}_r = \text{Cov}(\mathbf{v}_r | \hat{\mathbf{H}}_{\text{sr}}, \hat{\mathbf{H}}_{\text{rr}})$ and $\mathcal{Q}[l] \triangleq \{\mathbf{Q}_s[l], \mathbf{Q}_r[l]\}$. A similar expression is found for $\underline{I}_{\text{rd}}(\mathcal{Q}[l])$.

- We can also **upper-bound** the mutual information by ignoring the channel estimation error component.

Maximizing the Achievable-Rate Lower-Bound

- For full-duplex operation, the Source \rightarrow Destination rate is bottlenecked by the **smallest** of $\{I_{sr}, I_{rd}\}$.
- Therefore, our metric is

$$I_{\tau}(\mathcal{Q}) = \min \left\{ \underbrace{\sum_{l=1}^2 \tau[l] I_{sr}(\mathcal{Q}[l])}_{\triangleq I_{sr,\tau}(\mathcal{Q})}, \underbrace{\sum_{l=1}^2 \tau[l] I_{rd}(\mathcal{Q}[l])}_{\triangleq I_{rd,\tau}(\mathcal{Q})} \right\}$$

where $\mathcal{Q} \triangleq \{Q_s[1], Q_s[2], Q_r[1], Q_r[2]\}$.

- Our optimization problem becomes $\max_{\mathcal{Q}} I_{\tau}(\mathcal{Q})$ with **power and positivity constraints**

$$\mathcal{Q} \in \mathcal{Q}_{\tau} \triangleq \left\{ \begin{array}{l} \sum_{l=1}^2 \tau[l] \text{tr}(Q_s[l]) \leq 1 \quad , \quad Q_s[l] \geq \mathbf{0} \quad \forall l \in \{1, 2\} \\ \sum_{l=1}^2 \tau[l] \text{tr}(Q_r[l]) \leq 1 \quad , \quad Q_r[l] \geq \mathbf{0} \quad \forall l \in \{1, 2\} \end{array} \right\}.$$

Transmit Covariance Optimization

- We convert the maximin problem to a **weighted sum-rate** optimization problem

$$\max_{\zeta \in [0,1]} \max_{\mathcal{Q} \in \mathcal{Q}_\tau} (\zeta \underline{I}_{\text{sr},\tau}(\mathcal{Q}) + (1 - \zeta) \underline{I}_{\text{rd},\tau}(\mathcal{Q}))$$

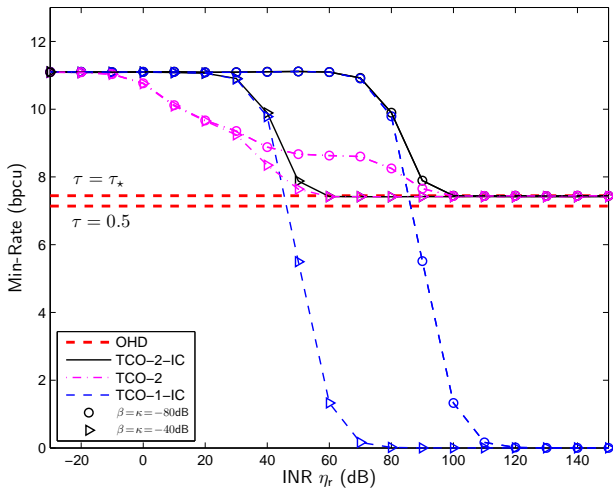
where we find ζ via **bisection search**.

- To maximize τ -weighted sum-rates $\underline{I}_{\text{sr},\tau}(\mathcal{Q})$ and $\underline{I}_{\text{rd},\tau}(\mathcal{Q})$, we have developed a **Gradient Projection** algorithm.
- The **projection** step is performing **waterfilling** over both spatial and temporal degrees of freedom.
- Finally we maximize with respect to the time-share τ using a **grid search**.

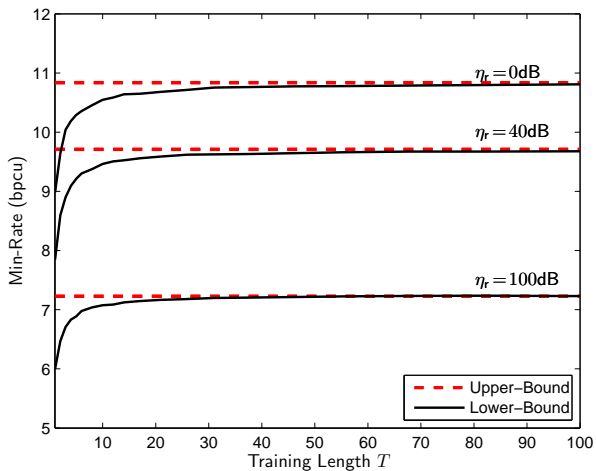
Numerical Results

- We will now show the **achievable-rate bounds** in the following plots:
 - versus INR η_r
 - versus training length
- In the plots, we show our proposed scheme as well as the following schemes:
 - **Half-duplex** with optimized covariance matrices and time-sharing parameter τ
 - Our proposed scheme **without performing interference cancellation**
 - Our proposed scheme using **only one period per data epoch**

Achievable-Rate Lower-Bound vs. INR η_r



Achievable-Rate Lower-Bound vs Training Length T



Conclusion

- We characterized the achievable-rate of MIMO decode-and-forward full-duplex relaying.
- We considered dynamic range limitations at the transmitter and receiver, as well as channel-estimation error from the training-based least-squares.
- Our solution required solving a non-convex optimization problem, for which we applied the projected gradient method.
- An analytic **approximation** that writes mutual information as an **explicit function of the SNRs, INRs, numbers of antennas, and dynamic-range parameters κ and β** was also derived (see paper).

Thanks!

Backup Slides

Gradient Projection Algorithm

We find the achievable-rate lower-bound via **Gradient Projection**:

for $k = 1, 2, 3, \dots$

$$\mathbf{P}_r^{(k)}[1] = \mathbf{Q}_r^{(k)}[1] + \mathbf{G}_r^{(k)}[1]$$

$$\mathbf{P}_r^{(k)}[2] = \mathbf{Q}_r^{(k)}[2] + \mathbf{G}_r^{(k)}[2]$$

$$(\tilde{\mathbf{Q}}_r^{(k)}[1], \tilde{\mathbf{Q}}_r^{(k)}[2]) = \mathcal{P}_{\mathbb{Q}_\tau}(\mathbf{P}_r^{(k)}[1], \mathbf{P}_r^{(k)}[2])$$

$$\mathbf{Q}_r^{(k+1)}[1] = \mathbf{Q}_r^{(k)}[1] + \gamma^{(k)}(\tilde{\mathbf{Q}}_r^{(k)}[1] - \mathbf{Q}_r^{(k)}[1])$$

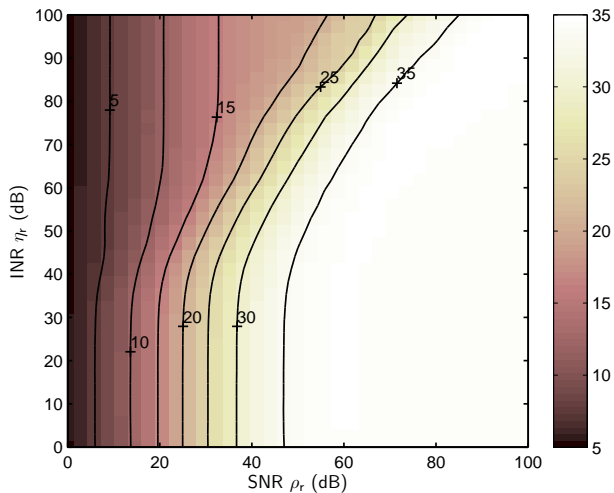
$$\mathbf{Q}_r^{(k+1)}[2] = \mathbf{Q}_r^{(k)}[2] + \gamma^{(k)}(\tilde{\mathbf{Q}}_r^{(k)}[2] - \mathbf{Q}_r^{(k)}[2])$$

⟨Similar repeated for $\mathbf{Q}_s[1]$ and $\mathbf{Q}_s[2]$ ⟩

end

where $\mathbf{G}_r^{(k)}[l]$ is the **gradient**, and $\mathcal{P}_{\mathbb{Q}_\tau}(\cdot)$ projects the period 1 and period 2 covariances onto the constraint set. $\gamma^{(k)}$ is chosen via the Armijo stepsize rule.

Achievable-Rate Lower-Bound Contour over SNR and INR



Analytic Approximation of Achievable Rate

- The complicated nature of the optimization problem motivates us to approximate its solution
- Making simplifying assumptions, we are able to find straightforward optimal transmit covariance matrices for both **full-duplex** and **half-duplex** operation.
- Our analytic approximate solution is simply the maximum of the **full-duplex** and **half-duplex** approximate solutions.

Analytic Approximation Contour over SNR and INR

